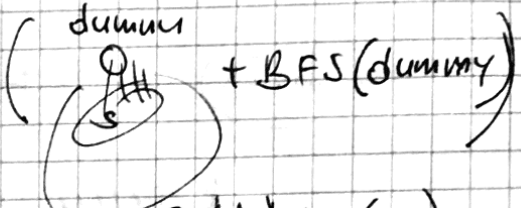


$$k = \sqrt{n}$$

defn  $N_k(u)$

(triangle inequality)  
implicitly used

1]  $S = \text{sample } |S| = d \frac{n}{k} \ln w$  nodes unif. at random  $d > 2$

2]  $w = \text{farthest node from } S$  

3] return  $\tilde{D} = \max_{v \in S \cup N_k(w)} \text{BFS}(v)$

Prop  $S$  is hitting set for  $N_k(u)$ ,  $u \in V$  w.h.p.

proof

$$P = \Pr(S \cap N_k(u) = \emptyset) = \binom{n-k}{n}^{|S|} = \left(1 - \frac{k}{n}\right)^{|S|} \approx e^{-\frac{k|S|}{n}} = e^{-\frac{k \cdot d \frac{n}{k} \ln w}{n}} = n^{-d}$$

#experiments  
stop BFS(w) after k nodes

Union bound:  $n \cdot n^{-d} = \frac{1}{n^{d-1}}$  at least one  $N_k(u)$  is not hit

$\Rightarrow$  prob.  $\geq 1 - \frac{1}{n^{d-1}}$  that  $S$  is hitting set  $\square$

w.h.p.  $d > 2$

Prop Let  $D = 3h + z$  where  $z \in \{0, 1, 2\}$

w.h.p.  $2h + z \leq \tilde{D} \leq D$   $z = 0, 1$   
 $z = 3$

proof Let  $a, b$  s.t.  $D = d(a, b)$  diametral nodes

2]  $d(w, S) \leq h \Rightarrow d(a, S) \leq h \Rightarrow \text{BFS}(x) \geq D - h = 2h + z$   
where  $x \in S$  is the closest point to  $a$   
or diameter  $> D$

(b)  $d(w, s) > h$  (i.e. every node in  $S$  is  $> h$  far from  $w$ )

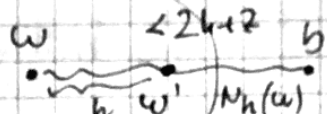
(b.1)  $\nexists$  BFS( $w$ )  $\geq 2h+2 \Rightarrow$  DONE

(b.2) BFS( $w$ )  $< 2h+2 \Rightarrow d(w, b) < 2h+2$

~~if~~  $S$  hits  $N_k(w)$  w.h.p.  $\Rightarrow$   $\bullet$

w.h.p.  $\exists x \in N_k(w)^{NS}$  s.t.  $d(w, x) > h \Rightarrow$

$N_k(w)$  contains <sup>at least</sup> all nodes  
at distance  $k$  from  $w$



w.h.p.  $\exists w'$  on the <sup>min</sup> path from  $w$  to  $b$  s.t.  $d(w, w') = h$

$$\Rightarrow d(w', b) < h+2 \Rightarrow d(w', s) \geq D - (h+2-1) = 2h+1$$

$$\Rightarrow \text{BFS}(w') \geq 2h+1 \quad \bullet$$