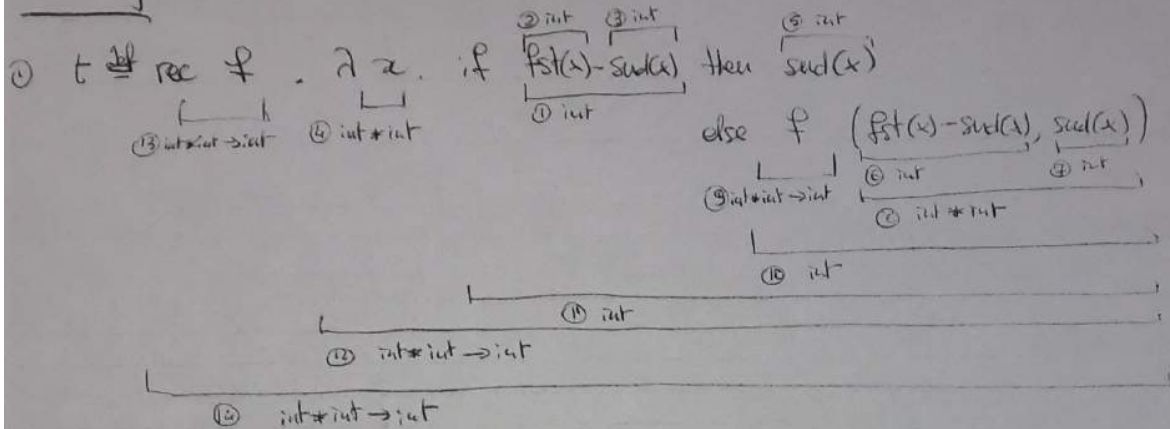


Ex3

Text



② $t(6,3) \rightarrow c \wedge t \rightarrow \lambda x. t' \quad t'[\frac{(6,3)}{x}] \rightarrow c$

$t' = \text{if } \text{fst}(6,3) - \text{snd}(6,3) \text{ then } \text{snd}(6,3)$
 $\text{else } t(\text{fst}(6,3) - \text{snd}(6,3), \text{snd}(6,3)) \rightarrow c$

$\text{fst}(6,3) - \text{snd}(6,3) \rightarrow m \quad m \neq 0 \quad t(\text{fst}(6,3) - \text{snd}(6,3), \text{snd}(6,3)) \rightarrow c$

$\text{fst}(6,3) \rightarrow m_1 \quad \text{snd}(6,3) \rightarrow m_2 \quad m_1 - m_2 \neq 0 \quad t(\text{fst}(6,3) - \text{snd}(6,3), \text{snd}(6,3)) \rightarrow c$

$m_1 = 6, m_2 = 3 \quad t(\frac{\text{fst}(6,3) - \text{snd}(6,3), \text{snd}(6,3)}{e}) \rightarrow c$

$t \rightarrow \lambda x. t'' \quad t''[\frac{e}{x}] \rightarrow c$

$t'' = \text{if } \text{fst}(e) - \text{snd}(e) \text{ then } \text{snd}(e)$
 $\text{else } t(\text{fst}(e) - \text{snd}(e), \text{snd}(e)) \rightarrow c$

$\text{fst}(e) - \text{snd}(e) \rightarrow 0 \quad \text{snd}(e) \rightarrow c$

$\text{fst}(e) \rightarrow m_1' \quad \text{snd}(e) \rightarrow m_2' \quad m_1' - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$e \rightarrow (t_0, t_1) \quad t_0 \rightarrow m_1' \quad \text{snd}(e) \rightarrow m_2' \quad m_1' - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$t_0 = \text{fst}(6,3) - \text{snd}(6,3), t_1 = \text{snd}(6,3) \quad \text{fst}(6,3) - \text{snd}(6,3) \rightarrow m_1' \quad \text{snd}(e) \rightarrow m_2' \quad m_1' - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$\text{snd}(e) \rightarrow m_2' \quad 3 - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$e \rightarrow (t_0', t_1') \quad t_1' \rightarrow m_2' \quad 3 - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$t_0' = \text{fst}(6,3) - \text{snd}(6,3), t_1' = \text{snd}(6,3) \quad \text{snd}(6,3) \rightarrow m_2' \quad 3 - m_2' = 0 \quad \text{snd}(e) \rightarrow c$

$3 - 3 = 0 \quad \text{snd}(e) \rightarrow c$

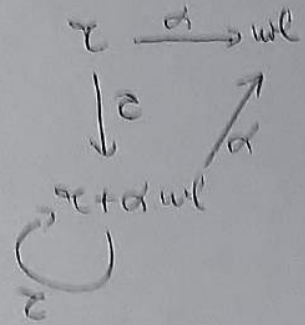
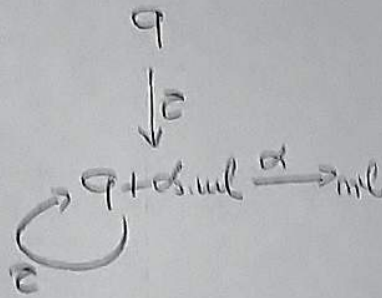
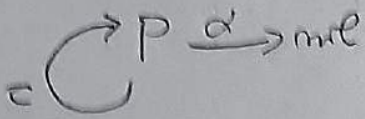
$c = 3 \quad \square$

Text

Es.5

Text

①



②

$$p \neq q$$

$$p \neq \alpha \text{ true}$$

$$q \neq \alpha \text{ true}$$

$$p \cong r$$

$R = \{ \{ p, r, r + \alpha.ml \}, \{ ml \} \}$ is a bisimulation

③

$$p \cong r$$

because $p \cong r$

$$p \cong q$$

$R = \{ \{ p, q, q + \alpha.ml \}, \{ ml \} \}$ is a weak bisimulation

